Box Progressions and Abelian Power Free Words

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The Setup

Given balls and boxes both enumerated with positive integers, we consider a sequential allocation of the balls into the boxes.

We fix an $\ell \geq 2$. Proceeding in increasing order of box labels, choose an arbitrary integer r between 1 and ℓ , and assign to each box the next r smallest unassigned balls.

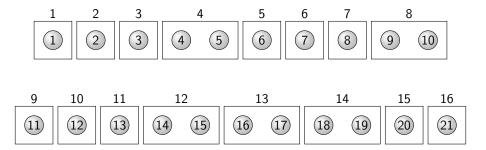
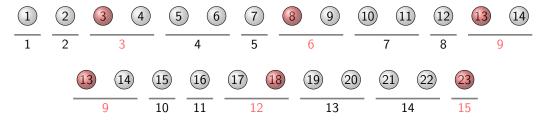


Figure 1: An example of a 2-ball-box distribution

- ▶ We focus on ball and box labels that simultaneously form a k-term arithmetic progression.
- \triangleright We call these balls and boxes k ball box progressions (k-BPs).



Balls: 3, 8, 13, 18, 23 Boxes: 3, 6, 9, 12, 15

Question

- We investigate the minimal number of balls needed to guarantee the existence of ball-box progressions
- ls there such a minimal number for every k?
- \triangleright Or are there infinitely long ball-box distributions with no k ball-box progressions?

- ightharpoonup l=2. First non-trivial case
- ▶ The problem seems related to Ramsey theory at first.

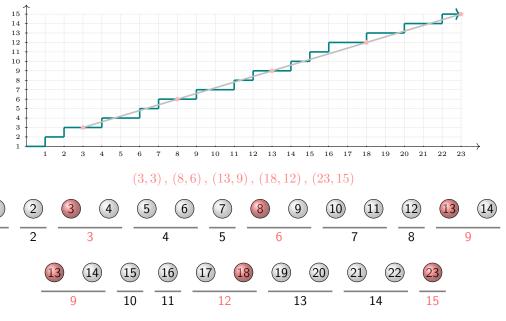
Plan Of The Talk

1 Introduction

2 Formalizing The Question

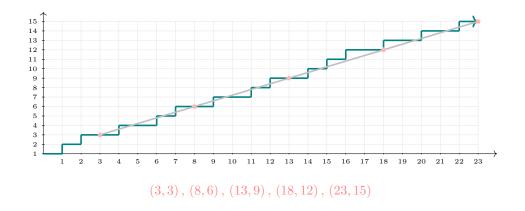
Combinatorics on Words

Equivalent Formalization- I



Balls: 3, 8, 13, 18, 23 Boxes: 3, 6, 9, 12, 15

Equivalent Formalization- I



- Indeed, we are looking for arithmetic progressions in $\mathbb{N} \times \mathbb{N}$ in the graph above.
- ► The analogy with Szemerédi is suggestive

Theorem (Multidimensional Szemerédi Theorem on \mathbb{N}^2)

Let $A \subset \mathbb{N}^2$ be a subset with positive upper density:

$$\overline{d}(A) = \limsup_{N \to \infty} \frac{|A \cap [1, N]^2|}{N^2} > 0.$$

Then, for every $k \in \mathbb{N}$, there exist $x \in \mathbb{N}^2$ and nonzero $v \in \mathbb{N}^2$ such that

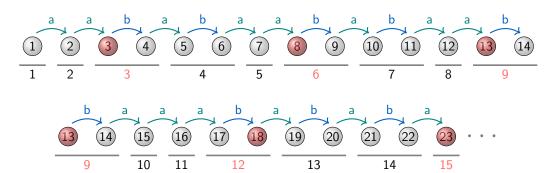
$$\{x, x+v, x+2v, \dots, x+(k-1)v\} \subset A.$$

- ▶ The density of integer points on the graph is $O\left(\frac{1}{N}\right)$. This is not a dense subset.
- Multidimensional Szemerédi is not applicable, another formalization?

Equivalent Formalization- II

We can encode this sequence into a word over $\Sigma := \{a, b\}$ via the following rules:

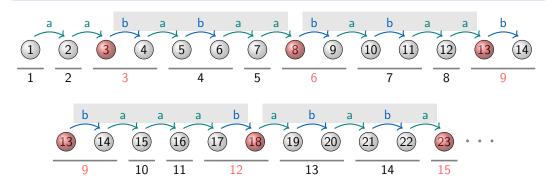
- ightharpoonup an a if the box changes
- ▶ a *b* if the same box



Equivalent Formalization- II

We can encode this sequence into a word over $\Sigma := \{a, b\}$ via the following rules:

- ightharpoonup an a if the box changes
- \triangleright a b if the same box



Combinatorics on Words

Repetition

A k repetition is composed of k consecutive equal factors. Examples:

- ► aaaa
- ► <u>aaabb</u> <u>aaabb</u> <u>aaabb</u> <u>aaabb</u>

Abelian Powe

A k abelian power is composed of k consecutive factors that are anagrams of each other. Examples:

- ► aaaa
- ► <u>aaabb</u> <u>aaabb</u> <u>aaabb</u> <u>aaabb</u>
- angle, angel
- ▶ <u>babaa babaa baaab ababa</u>

Ball Box Progressions and Abelian Powers

Proposition

Let $k \geq 2$ be a positive integer. Let $\mathcal D$ a ball-box distribution, the corresponding word $w_{\mathcal D}$ is abelian (k-1)-power free if and only if there is no k ball-box progression in $\mathcal D$.

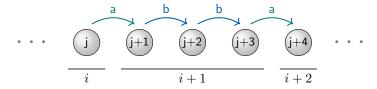
being k ball-box progression free \longleftrightarrow being k-1 abelian power free

Some previous work on abelian power-free words:

- F. Dekking, Strongly non-repetitive sequences and progression-free sets, Journal of Combinatorial Theory 27 (1979)
- A. Carpi, On abelian power-free morphisms, Internat. J. Algebra Comput. 3 (1993) 151–167.

Our Constraint On The Word

at most 2 balls in a box \longleftrightarrow no bb as a factor



1 Introduction

Pormalizing The Question

3 Combinatorics on Words

Some Notions on Words-I

- Let $\Sigma = \{a, b\}$ be an alphabet. An element of Σ^* is called a word.
- ▶ The empty word is denoted by ϵ .
- Given a word

$$v = uw$$
,

u is called a prefix.

We set

$$\operatorname{Pref}(v) = \{u \in \Sigma^* : u \text{ is a prefix of } v\}.$$

ightharpoonup The Parikh vector of v denoted by

$$\psi(v) = (v_1, v_2) \in \mathbb{N}^{|\Sigma|}$$

where v_1 and v_2 is the number of appearances of a and b respectively.

Some Notions on Words-II

► A morphism is a map

$$h: \Sigma^* \to \Sigma^*$$

 $h(vw) \mapsto h(v)h(w).$

► The frequency matrix of a morphism is

$$\begin{pmatrix} \psi(h(a)) \\ \psi(h(b)) \end{pmatrix}.$$

Finally, we define

$$\operatorname{Pref}(h) := \bigcup_{\sigma \in \Sigma} \operatorname{Pref}(h(\sigma)).$$

▶ h is called abelian k-power-free if, for every abelian k-power free word $w \in \Sigma^*$, the image h(w) is also ableian k-power free.

Example

 $lackbox{ Consider the morphism } h:\Sigma^* o\Sigma^* \mbox{ given by }$

$$h(a) = aaaba$$
$$h(b) = bab.$$

The frequency matrix of h is

$$M = \begin{pmatrix} 4 & 1 \\ 1 & 2 \end{pmatrix}$$

The determinant of the matrix is 7.

$$\operatorname{Pref}(h) = \{\epsilon, a, aa, aaa, aaab, b, ba\}$$

Questior

ls h a k abelian power free morphism for some k?

Some Observation on h

▶ Define the group generated by $\psi(h(a)) = (4,1)$ and $\psi(h(b)) = (1,2)$.

$$G_h = \langle \psi(h(\sigma)) : \sigma \in \Sigma \rangle \subseteq \mathbb{Z}^{\Sigma}.$$

$$\operatorname{Pref}(h) = \{ \epsilon, a, aa, aaa, aaab, b, ba \}$$

$$\mathcal{P} := \psi(\operatorname{Pref}(h)) = \{ (0, 0), (1, 0), (2, 0), (3, 0), (3, 1), (0, 1), (1, 1) \}.$$

▶ We observe that (0,1) and (1,2) are in the same coset in \mathbb{Z}^{Σ}/G_h :

$$(3,0) - (0,1) = (4,1) - (1,2).$$

$$\overline{\mathcal{P}} = \{ \overline{(0,0)}, \overline{(1,0)}, \overline{(2,0)}, \overline{(3,0)}, \overline{(1,1)}, \overline{(3,1)} \}.$$

Arithmetic Rank (a-rk)

The arithmetic rank of a set A is defined as

 $\operatorname{a-rk}(A) := \text{length of the longest arithmetic progression in } A$

ightharpoonup Example: $\operatorname{a-rk}(\psi(\operatorname{Pref}(h))) = 6$

Theorem (Eyidoğan, Göral, T.)

Given an integer $n \geq 4$, an alphabet Σ of size at least 2 and a morphism

$$h: \Sigma^* \to \Sigma^*$$
.

Let

$$G_h = \langle \psi(h(\sigma)) : \sigma \in \Sigma \rangle \subseteq \mathbb{Z}^{\Sigma}.$$

Finally we set

$$\mathcal{P}=\psi(\mathrm{Pref}\,(h)) \text{ and } \overline{\mathcal{P}}=\{\overline{v}:v\in\mathcal{P}\}.$$

Suppose the following conditions are satisfied:

- 1. The frequency matrix of h is non-singular,
- 2. $\operatorname{a-rk}(\overline{\mathcal{P}}) \leq n$,
- 3. For $v \in \mathcal{P} \setminus (\psi(h(\Sigma)) \cup \{0\})$, the equivalence class of v contains at most 2 elements. Moreover if v and w are in the same class, then

$$v - w = \psi(h(p)) - \psi(h(q))$$

for some $p,q \in \Sigma$ where $v \in \psi(\operatorname{Pref}(h(p)))$ and $w \in \psi(\operatorname{Pref}(h(q)))$.

Then, h is an abelian n-power free morphism.

Corollary

There exists an infinite word Ω over a binary alphabet $\Sigma = \{a, b\}$ such that

- $\blacktriangleright bb$ is not a factor of Ω (there are no two consecutive b's in Ω),
- $ightharpoonup \Omega$ is abelian 6-power free.
- ▶ We consider the morphism

$$h(a) = aaaba \ h(b) = bab.$$
 with non-singular frequency matrix $\begin{pmatrix} 4 & 1 \ 1 & 2 \end{pmatrix}$.

- $ightharpoonspice a-rk(\overline{\mathcal{P}})=6$
- ▶ We have only two elements in the same equivalence class:

$$(3,0) - (0,1) = (4,1) - (1,2).$$

and

$$(1,2) - (0,1) = (4,1) - (3,0) = \psi(h(a)) - \psi(h(b)).$$

In h is an abelian power-free morphism, hence $\Omega:=h^\omega(a)$ is an infinite abelian 6 power-free word.

Answer

- $ightharpoonup \Omega$ admits no 6 abelian powers
- Among the ball-box distributions where we allow at most 2 balls in a box, there is an infinitely long ball-box distribution corresponding to Ω which admits no 7 ball box progressions.

